# Concurrent Cube-and-Conquer 

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## Introduction

- lookahead (LA)
- recursively split instance: binary search tree
- good for small, hard problems
- conflict-driven clause learning (CDCL)
- learn implied clauses: less systematic search
- best for large, "easy" industrial instances
- cube-and-conquer (CC)
- partition using LA into thousands or millions of subproblems
- solve subproblems in parallel using CDCL


## Cube-and-conquer



## Cube-and-conquer



## Old CC cutoff heuristic

$d\left(c_{\text {id }}\right):=\left|\varphi_{\text {dec }}\right|^{2} \cdot\left(\left|\varphi_{\text {dec }}\right|+\left|\varphi_{\text {imp }}\right|\right)$
$d\left(c_{\text {id }}\right)>$ threshold $\rightarrow$ cut off
Dynamic threshold:

- LA refutes cube $\rightarrow$ decrease
- increase gradually


## Motivation

Limitations of cube-and-conquer (CC):

- partitioning not ideal
- lookahead not always effective

Proposed solutions:

- run CDCL and LA concurrently in partitioning phase
- predict unsuitable instances


## Concurrent cube-and-conquer

Solve by adding CDCL to cube phase:

- run LA and CDCL concurrently
- add decisions by LA as assumptions to CDCL
- use existing solvers (March_rw and MiniSAT 2.2)


## Concurrent cube-and-conquer


(Cx) refuted by LA
(Cx) refuted by CDCL

## Concurrent cube-and-conquer


(Cx) refuted by LA
(Cx) refuted by CDCL

## Concurrent cube-and-conquer


(cx) current
(Cx) refuted by LA
(Cx) refuted by CDCL

## Concurrent cube-and-conquer


(c) cube
(cx) current
(Cx) refuted by LA
(Cx) refuted by CDCL

## Concurrent cube-and-conquer


(cx) current
(cx refuted by LA
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## Concurrent cube-and-conquer


(Cx) current
(Cx) refuted by LA
(Cx) refuted by CDCL

## Concurrent cube-and-conquer

(c) cube

( $C_{x}$ current
(Cx) refuted by LA
(Cx) refuted by CDCL

## Concurrent cube-and-conquer


$c_{x}$ current
(cx refuted by LA
(Cx) refuted by CDCL

## Concurrent cube-and-conquer

(c) cube

( $C_{x}$ current
(Cx) refuted by LA
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## Concurrent cube-and-conquer


$C_{X}$ cube
$c_{x}$ current

(cx) refuted by LA
(C) refuted by CDCL

## Concurrent cube-and-conquer


$C_{X}$ cube
$c_{x}$ current

(Cx) refuted by LA
(Cx) refuted by CDCL

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$c_{x}$ current

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(cx) refuted by LA
(cx refuted by CDCL

## Concurrent cube-and-conquer


(Cx) cube
(Cx) current

unsat
$\left\langle c_{4}\right\rangle$

## Concurrent cube-and-conquer


(cx) refuted by LA
(cx refuted by CDCL

## Concurrent cube-and-conquer


(cx) refuted by LA
(cx refuted by CDCL

## Concurrent cube-and-conquer


(cx) refuted by LA
(cx) refuted by CDCL

Use information from the CDCL solver in CCC's cutoff heuristic
Like in CC, $d\left(c_{\mathrm{id}}\right)>$ threshold $\rightarrow$ cut off
But now:

- CDCL refutes cube $\rightarrow$ decrease threshold
- LA refutes cube $\rightarrow$ increase threshold (vs. decrease in CC)


## Motivation

Limitations of cube-and-conquer (CC):

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Proposed solutions:

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- predict unsuitable instances


## CCC without prediction




Run time of MiniSAT (vertical) and $\mathrm{CCC}_{\text {mini }}$ (horizontal) on SAT 2009 and 2011 instances.
Application (left) and crafted (right) instances.

## Predicting effectiveness of (C)CC

Lookahead effective if

- lookahead refutes cubes before CDCL
- limited number of right branches

Use pure CDCL if CCC seems ineffective after 5 seconds


## Predicting effectiveness of (C)CC




Run time of MiniSAT (vertical) and $\mathrm{CCC}_{\text {mini }}$ (horizontal) on instances selected (■) and not selected ( $\boldsymbol{\Delta}$ ) by the predictor. Application (left) and crafted (right) instances.

## Results on filtered crafted instances



## Results on filtered application instances



## Conclusion

CCC solves CC's limitations

- CCC uses CDCL to find a better cutoff point
- CCC switches to pure CDCL if partitioning performs poorly

And is

- often faster than CDCL, LA, and CC
- natural to parallelize
- and ready to compete in the SAT Challenge 2012


# Concurrent Cube-and-Conquer 

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